

Complete Identification of Generators and Check Elements of Zero Divisor Codes over Cyclic Group Rings

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Abstract

Let F_q be a finite field with $q = p^n$ where p is a prime and C_m be a cyclic group of order m with generator g . Let u be a zero divisor in a cyclic group ring $F_q C_m$ having W as a F_q -submodule. Then $C = Wu$ is called a zero divisor code over $F_q C_m$ with generator u . To implement zero divisor codes in real world, each C needs a generator u as well as explicitly identifying its respective check element which is a principle zero divisor partner of u . In this paper, an algebraic study on zero divisors of $F_q C_m$ is done by viewing each $u = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q C_m$ as the polynomial $u(g) = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q[g]$. It is found that $\gcd(u(g), g^m - 1)$ plays a vital role in extracting important properties of zero divisors in $F_q C_m$. Using the obtained results, the set of all zero divisors in $F_q C_m$ or all generators of zero divisor codes over $F_q C_m$ is completely identified with its cardinality stated explicitly. This paper ends with an explicit identification of a check element for each zero divisor code over $F_q C_m$.

Key words and phrases: Cyclic Group Ring, Principal Zero Divisor Partner, Zero Divisor Code, Check Element

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1 Introduction

Let F_q be a finite field with $q = p^n$ where p is a prime and $G = \{g_0 = 1, g_1, \dots, g_{m-1}\}$ be an abelian group of order m . Then $F_qG = \{\sum_{g \in G} \alpha_g g \mid \alpha_g \in F_q\}$ is called a group ring of G over F_q , which has algebraic structures of both commutative ring and F_q -module. In addition, let $M_{m \times m}(F_q)$ be the ring of $m \times m$ matrices over F_q . With respect to the ordered listing of elements in G , there exists an injective homomorphism $\sigma_G : F_qG \rightarrow M_{m \times m}(F_q)[1]$. For each $a \in F_qG$, the $\sigma_G(a)$ is called the group ring matrix of a (with respect to the ordered listing G). Clearly, different ordered listings of elements in G give different group ring matrices of a . However, all group ring matrices of a share the same rank called the rank of a , denoted as $rank(a)$, that is $rank(a) = rank(\sigma_G(a))$.

Let W be a F_q -submodule of F_qG and $u \in F_qG$ be a zero divisor. Then for all $x \in W$, $f_{W,u} : W \rightarrow F_qG$ such that $(x)f_{W,u} = xu$ is an encoding linear map for a code $C = Im(f_{W,u}) = \{xu \mid x \in W\} = Wu$. Note that the respective code is a F_q -submodule of F_qG and it is named zero divisor code with a generator u relative to the F_q -submodule W [2].

Recall that a code with larger minimal distance can detect or correct more errors [5]. For non-isomorphic groups G of order $m \in \mathbb{N}$, with respects to a F_q -submodule W , different generators give different minimal distances for the zero divisor codes over F_qG . Hence, it is a necessity to study algebraic methods to completely identify all generators for zero divisor codes over F_qG for non-isomorphic groups G . Idempotents form a special class of zero divisors in F_qG . Recently, a complete algebraic identification of all idempotents in F_qG and some interesting algebraic results on zero divisor codes with idempotent generators have been obtained [6, 7, 8].

On the other hand, in the decoding of $C = Wu$, a check element of C is an element $v \in F_qG$ such that $C = \{a \in F_qG \mid av = 0\}$ [2]. In other words, for a respective received message $w \in F_qG$ if $wv \neq 0$, then clearly w is received with errors. Let $u' \in F_qG$ be a zero divisor partner of u . Then $uu' = 0$ or $\sigma_G(u)\sigma_G(u') = 0$ implies the row space of $\sigma_G(u)$ is a subspace of the null space of $\sigma_G(u')$. Thus $rank(\sigma_G(u')) \leq |G| - rank(\sigma_G(u))$ or $rank(u') \leq |G| - rank(u)$. Note that $rank(u') < |G| - rank(u)$ if and only if $C \subsetneq \{a \in F_qG \mid au' = 0\}$. Hence, a $u' \in F_qG$ such that $rank(u') = |G| - rank(u)$ is a natural check element of C and thus it is specifically denoted as v , called principal zero divisor partner of u [2].

From the algebraic point of view, it is interesting to investigate the family of groups that assures the existence of a principal zero divisor partner for

each zero divisor u in the respective group ring [3, 4]. Till now, it is still a great challenge to convert the obtained algebraic results into an algebraic algorithm to explicitly identify a check element for each zero divisor code over the respective group ring.

Let $C_m = \langle g \rangle = \{g^0 = 1, g, g^2, \dots, g^{m-1}\}$ be a cyclic group of order m with generator g and thus $F_q C_m = \{\sum_{i=0}^{m-1} \alpha_i g^i | \alpha_i \in F_q\}$. Let $F_q[g]$ be the ring of polynomials over F_q with indeterminate g . Note that

$$\begin{aligned} F_q[g]/\langle g^m - 1 \rangle &= \left\{ \left(\sum_{i=0}^{m-1} \alpha_i g^i \right) + \langle g^m - 1 \rangle \middle| \alpha_i \in F_q \right\} \\ &= \left\{ \sum_{i=0}^{m-1} [(\alpha_i + \langle g^m - 1 \rangle) (g^i + \langle g^m - 1 \rangle)] \middle| \alpha_i \in F_q \right\}. \end{aligned}$$

As $F_q \cong \bar{F}_q = \{\alpha + \langle g^m - 1 \rangle | \alpha \in F_q\}$ and $C_m \cong G = \langle g + \langle g^m - 1 \rangle \rangle$, it can be seen that $F_q C_m$ is isomorphic to $F_q[g]/\langle g^m - 1 \rangle$ [9]. Throughout this paper, let $a(g) = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q[g]$ be the polynomial form of a group ring element $a = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q C_m$. However, any results obtained through the polynomial form will need to be reduced modulo $g^m - 1$ before it is applied into $F_q C_m$. In [3], Hurley proved that if $u(g) | (g^m - 1)$, then $u' \in F_q C_m$ such that $u(g)u'(g) = g^m - 1$ is a principal zero divisor partner of u .

In this paper, an algebraic method to explicitly identify a check element of $C = Wu$ (or a principal zero divisor partner of $u \in F_q C_m$) regardless of whether $u(g)$ is a divisor of $g^m - 1$ is discussed. Before that, for the completeness of our study, an algebraic method to completely identify all generators for zero divisor codes over $F_q C_m$ or all zero divisors in $F_q C_m$ is introduced.

2 Identification of Generators For Zero Divisor Codes Over $F_q C_m$

As the complete identification of generators for zero divisor codes over $F_q C_m$ implies the complete identification of zero divisors in $F_q C_m$ and vice-versa, a practical method is firstly developed to sort out zero divisors among non-zero elements in $F_q C_m$.

Let a be a non-zero element in $F_q C_m$. We show that the determination of whether a is a unit or a zero divisor can be done by using the greatest

common divisor of $a(g)$ and $g^m - 1$ that is $\gcd(a(g), g^m - 1)$. Before that, we need the following result:

Lemma 2.1. *Let $a \in F_q C_m$. Then $\gcd(a(g), g^m - 1) = 1$ if and only if there exist $\beta, \eta \in F_q C_m$ such that $\beta(g)a(g) + \eta(g)(g^m - 1) = 1$.*

Proof. Assume that $\gcd(a(g), g^m - 1) = 1$. Then from [5], it is known that there exist $\beta(g) = \sum_{i=0}^{m-1} \beta_i g^i, \eta(g) = \sum_{i=0}^{m-1} \eta_i g^i \in F_q[g]$, equivalently $\beta, \eta \in F_q C_m$, such that $\beta(g)a(g) + \eta(g)(g^m - 1) = 1$. Conversely, assume that there exist $\beta, \eta \in F_q C_m$, such that $\beta(g)a(g) + \eta(g)(g^m - 1) = 1$. Let $\gcd(a(g), g^m - 1) = d(g)$. Note that since $d(g)|a(g)$ and $d(g)|g^m - 1$, then it follows that $d(g)|(\beta(g)a(g) + \eta(g)(g^m - 1))$ which implies $d(g)|1$. Hence, $d(g) = 1$. \square

Theorem 2.2. *Let $a \in F_q C_m - \{0\}$. Then a is a unit if and only if $\gcd(a(g), g^m - 1) = 1$. Otherwise, a is a zero divisor.*

Proof. By Lemma 2.1, $\gcd(a(g), g^m - 1) = 1$ if and only if there exist $\beta, \eta \in F_q C_m$ such that $a(g)\beta(g) + (g^m - 1)\eta(g) = 1$ if and only if $a(g)\beta(g) = 1 - (g^m - 1)\eta(g) \equiv 1 \pmod{g^m - 1}$ if and only if $a\beta = 1$ (recall that $F_q C_m \cong F_q[g]/\langle g^m - 1 \rangle$) if and only if a is a unit. As each non-zero element in $F_q C_m$ is either a unit or a zero divisor [1], the result of a is a zero divisor if and only if $\gcd(a(g), g^m - 1) \neq 1$ followed. \square

Using Theorem 2.2, note that $\gcd(a(g), g^m - 1) \neq 1$ if and only if $1 \leq \deg(a(g)) \leq m - 1$ and $a(g)$ is a multiple of some irreducible factors of $g^m - 1$. Thus it is clear that the initial step to detect all zero divisors of $F_q C_m$ is by identifying the complete factorization of $g^m - 1$ into product of irreducible polynomials over F_q .

Let $m \in \mathbb{N}$ such that $\gcd(m, q) = 1$. Recall that for each $0 \leq j \leq m - 1$, the cyclotomic coset of q modulo m containing j is $C_j = \{jq^i \pmod{m} | i = 0, 1, 2, \dots\}$ [5]. The following is a result that can be used to identify all irreducible factors of $g^m - 1$ over F_q for the case of $\gcd(m, q) = 1$.

Theorem 2.3. [5] *Let $m \in \mathbb{N} - 1$ such that $\gcd(m, q) = 1$. Suppose that n is a positive integer satisfying $m|(q^n - 1)$. Let α be a primitive element of F_{q^n} and $M_j(g) = \prod_{i \in C_j} (g - \alpha^i)$ be the minimal polynomial of α^j with respect to F_q . Let $\{s_1, \dots, s_t\}$ be a complete set of representatives of cyclotomic cosets of q modulo m . Then $g^m - 1$ has a total of t distinct monic irreducible factors over F_q where $g^m - 1 = \prod_{i=1}^t M_{((q^n - 1)s_i/m)}(g)$.*

Consider the case of $\gcd(m, q) \neq 1$. Recall that for $\alpha_1, \alpha_2 \in F_q$, $(\alpha_1 + \alpha_2)^{p^\vartheta} = \alpha_1^{p^\vartheta} + \alpha_2^{p^\vartheta}$ for every $\vartheta \in \mathbb{N}$ [3]. Note that as a positive integer and

under the condition of $\gcd(m, q) \neq 1$, $m = m'p^\mu$ for some $\mu, m' \in \mathbb{N}$ with $\gcd(m', p) = 1$. Then $g^m - 1 = g^{m'p^\mu} - 1^{p^\mu} = (g^{m'} - 1)^{p^\mu}$.

Case 1. $m' = 1$ that is $m = p^\mu$ for some $\mu \in \mathbb{N}$. Then $g^m - 1 = g^{p^\mu} - 1^{p^\mu} = (g - 1)^{p^\mu}$

Case 2. $m' > 1$. Let n be a positive integer satisfying $m'|(q^n - 1)$. Then by Theorem 2.3, $g^{m'} - 1 = \prod_{i=1}^t M_{((q^n-1)s_i/m')}(g)$ and thus

$$\begin{aligned} g^m - 1 &= g^{m'p^\mu} - 1^{p^\mu} \\ &= (g^{m'} - 1)^{p^\mu} \\ &= \left(\prod_{i=1}^t M_{((q^n-1)s_i/m')}(g) \right)^{p^\mu} \\ &= \prod_{i=1}^t (M_{((q^n-1)s_i/m')}(g))^{p^\mu}. \end{aligned}$$

Hence, we have proved the following result:

Theorem 2.4. Suppose that m is a positive integer such that $\gcd(m, q) \neq 1$. Let $m = m'p^\mu$ for some $\mu, m' \in \mathbb{N}$ with $\gcd(m', p) = 1$. Let n be a positive integer satisfying $m'|(q^n - 1)$. Then

- (i) $g^m - 1 = (g - 1)^{p^\mu}$ if $m' = 1$;
- (ii) $g^m - 1 = \prod_{i=1}^t (M_{((q^n-1)s_i/m')}(g))^{p^\mu}$ if $m' > 1$.

From now on, let $m = m'p^\mu$ for some non-negative integer μ . Then $\gcd(m, q) = 1$ if $\mu = 0$ or $\gcd(m, q) \neq 1$ if $\mu \geq 1$. Recall from Theorem 2.4 that $g^m - 1 = \prod_{j=1}^t (M_j(g))^{p^\mu}$ with each $M_j(g)$ is a monic irreducible polynomial over F_q . For each $j \in \{1, 2, \dots, t\}$, define

$$\Upsilon_j := \left\{ b(g)M_j(g) \left| b(g) = \sum_{k=0}^{m-1-\deg(M_j(g))} \epsilon_k g^k, \epsilon_k \in F_q \right. \right\}. \quad (2.1)$$

Note that for each $1 < s \leq p^\mu$,

$$\left\{ b(g)M_j(g)^s \mid b(g) = \sum_{k=0}^{m-1-s \cdot \deg(M_j(g))} \epsilon_k g^k, \epsilon_k \in F_q \right\} \\ \subset \left\{ b(g)M_j(g) \mid b(g) = \sum_{k=0}^{m-1-\deg(M_j(g))} \epsilon_k g^k, \epsilon_k \in F_q \right\}.$$

Hence, the set of all zero divisors in $F_q C_m$ is $\bigcup_{j=1}^t \Upsilon_j - \{0\}$.

Theorem 2.5. *Let $g^m - 1 = \prod_{j=1}^t (M_j(g))^{p^\mu}$ where p is prime, $\mu \geq 0$ and each $M_j(g)$ is a monic irreducible polynomial over F_q . For each $j \in \{1, 2, \dots, t\}$, let Υ_j be defined as in Equation (2.1). Then $Z_{F_q C_m} = \bigcup_{j=1}^t \Upsilon_j - \{0\}$ where $Z_{F_q C_m}$ is the set of all zero divisors in $F_q C_m$.*

Corollary 2.6. *The set of all generators for zero divisor codes over $F_q C_m$ is $Z_{F_q C_m}$.*

Next, the cardinality of $Z_{F_q C_m}$ is studied. By Theorem 2.5, $|Z_{F_q C_m}| = \left| \bigcup_{j=1}^t \Upsilon_j - \{0\} \right|$. As $\Upsilon_{j_1} \cap \Upsilon_{j_2}$ is not necessary an empty set for all $j_1, j_2 \in \{1, 2, \dots, t\}$, $|Z_{F_q C_m}| \leq \sum_{j=1}^t |\Upsilon_j| - 1$. Hence, inclusion-exclusion principle is used to find $|Z_{F_q C_m}|$ by expanding the notation from Equation (2.1). For each $\lambda \in \{1, 2, \dots, t\}$, define

$$\Upsilon_{j_1, \dots, j_\lambda} := \left\{ b(g) (M_{j_1}(g) \cdots M_{j_\lambda}(g)) \mid b(g) = \sum_{k=0}^{m-1-\sum_{r=1}^\lambda \deg(M_{j_r}(g))} \epsilon_k g^k, \epsilon_k \in F_q \right\}. \tag{2.2}$$

For each $\lambda \in \{2, \dots, t\}$, it is obvious that $\Upsilon_{j_1, \dots, j_\lambda} \subseteq \bigcap_{k=1}^\lambda \Upsilon_{j_k}$. Conversely, as each $M_{j_k}(g)$ is irreducible, then if $a(g) \in \bigcap_{k=1}^\lambda \Upsilon_{j_k}$, $a(g) = b(g) (M_{j_1}(g) \cdots M_{j_\lambda}(g))$ for some $b(g) = \sum_{k=0}^{m-1-\sum_{r=1}^\lambda \deg(M_{j_r}(g))} \epsilon_k g^k$. Thus for each $\lambda \in \{2, \dots, t\}$, $\bigcap_{k=1}^\lambda \Upsilon_{j_k} \subseteq \Upsilon_{j_1, \dots, j_\lambda}$. Therefore, $\Upsilon_{j_1, \dots, j_\lambda} = \bigcap_{k=1}^\lambda \Upsilon_{j_k}$ for each $\lambda \in \{2, \dots, t\}$.

In particular, consider the case of $\mu = 0$. Let $a(g) \in \bigcap_{j=1}^t \Upsilon_j$. Then as each $M_{j_k}(g)$ is irreducible and $\deg(a(g)) \leq m-1$, $a(g) = b(g) (M_1(g) \cdots M_t(g))$ implies that $b(g) = 0$. Hence, the following result is obtained:

Lemma 2.7. *Let $g^m - 1 = \prod_{j=1}^t (M_j(g))^{p^\mu}$ where p is prime, $\mu \geq 0$ and each $M_j(g)$ is a monic irreducible polynomial over F_q . Let Υ_j be defined*

as in Equation (2.1) for each $j \in \{1, 2, \dots, t\}$ and $\Upsilon_{j_1, \dots, j_\lambda}$ be defined as in Equation (2.2) for each $\lambda \in \{1, 2, \dots, t\}$. Then $\bigcap_{k=1}^\lambda \Upsilon_{j_k} = \Upsilon_{j_1, \dots, j_\lambda}$ for each $\lambda \in \{2, \dots, t\}$. In addition, if $\mu = 0$, then $\bigcap_{j=1}^t \Upsilon_j = \{0\}$.

The next result gives the value of $|Z_{F_q C_m}|$.

Theorem 2.8. Let $g^m - 1 = \prod_{j=1}^t (M_j(g))^{p^\mu}$ where p is prime, $\mu \geq 0$ and each $M_j(g)$ is a monic irreducible polynomial over F_q . Then

$$|Z_{F_q C_m}| = \left(\sum_{\lambda=0}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - 1.$$

Proof. By Theorem 2.5, $|Z_{F_q C_m}| = |\bigcup_{j=1}^t \Upsilon_j - \{0\}|$. Then by inclusion-exclusion principle and Lemma 2.7, we have

$$\begin{aligned} \left| \bigcup_{j=1}^t \Upsilon_j \right| &= \sum_{j=1}^t |\Upsilon_j| - \sum_{1 \leq j_1 < j_2 \leq t} \left| \bigcap_{k=1}^2 \Upsilon_{j_k} \right| + \sum_{1 \leq j_1 < j_2 < j_3 \leq t} \left| \bigcap_{k=1}^3 \Upsilon_{j_k} \right| - \dots + (-1)^{t+1} \left| \bigcap_{j=1}^t \Upsilon_j \right| \\ &= \sum_{j=1}^t |\Upsilon_j| - \sum_{1 \leq j_1 < j_2 \leq t} |\Upsilon_{j_1, j_2}| + \sum_{1 \leq j_1 < j_2 < j_3 \leq t} |\Upsilon_{j_1, j_2, j_3}| - \dots + (-1)^{t+1} |\Upsilon_{1, 2, \dots, t}| \\ &= \sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} |\Upsilon_{j_1, \dots, j_\lambda}| \right) \\ &= \sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right). \end{aligned}$$

Thus,

$$\begin{aligned} |Z_{F_q C_m}| &= \left| \bigcup_{j=1}^t \Upsilon_j - \{0\} \right| \\ &= \left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - 1. \end{aligned}$$

□

For the case of $\gcd(m, q) \neq 1$, recall that $m = m'p^\mu$ for some $\mu \geq 1$ with $\gcd(m', p) = 1$. The following result gives the relation between $|Z_{F_q C_m}|$ and $|Z_{F_q C_{m'}}|$:

Corollary 2.9. Let $m = m'p^\mu \in \mathbb{N}$ such that p is prime and $\mu > 0$. Then $|Z_{F_q C_m}| = q^{(p^\mu - 1)m'} (|Z_{F_q C_{m'}}| + 1) - 1$.

Proof. Recall from Theorem 2.8 that

$$|Z_{F_q C_m}| = \left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - 1 \text{ and}$$

$$|Z_{F_q C_{m'}}| = \left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m' - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - 1.$$

Then

$$\begin{aligned} |Z_{F_q C_m}| &= \left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{p^\mu m' - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - 1 \\ &= \left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{(p^\mu - 1)m'} q^{m' - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) \right) - \\ &\quad q^{(p^\mu - 1)m'} + q^{(p^\mu - 1)m'} - 1 \\ &= q^{(p^\mu - 1)m'} \left(\left(\sum_{\lambda=1}^t (-1)^{\lambda+1} \left(\sum_{1 \leq j_1 < \dots < j_\lambda \leq t} q^{m' - \sum_{r=1}^\lambda \deg(M_{j_r}(g))} \right) - 1 \right) + 1 \right) - 1 \\ &= q^{(p^\mu - 1)m'} (|Z_{F_q C_{m'}}| + 1) - 1. \end{aligned}$$

□

Example 2.10. Consider $F_2 C_3$ with $\gcd(3, 2) = 1$. By Theorem 2.3, $g^3 - 1 = (g^2 + g + 1)(g + 1)$ with $M_1(g) = g^2 + g + 1$ and $M_2(g) = g + 1$ are all monic irreducible polynomials over F_2 . Thus $\Upsilon_1 = \{g^2 + g + 1, 0\}$ and $\Upsilon_2 = \{g + 1, g^2 + 1, g^2 + g, 0\}$ as defined in Equation (2.1). Let $\chi_1 = \Upsilon_1 - \{0\} = \{g^2 + g + 1\}$ and $\chi_2 = \Upsilon_2 - \Upsilon_1 = \{g + 1, g^2 + 1, g^2 + g\}$. Note that $\chi_1 \cap \chi_2 = \emptyset$ and $\bigcup_{j=1}^2 \chi_j = \bigcup_{j=1}^2 \Upsilon_j - \{0\}$ that is $\{\chi_1, \chi_2\}$ forms a partition of $Z_{F_2 C_3}$ by Theorem 2.5. Thus $|Z_{F_2 C_3}| = 4$ with $Z_{F_2 C_3} = \{g + 1, g^2 + 1, g^2 + g, g^2 + g + 1\}$. Note that this result is consistent with the value of $|Z_{F_2 C_3}|$ obtained using Theorem 2.8 that is

$$\begin{aligned} |Z_{F_2 C_3}| &= (2^{3 - \deg(M_1(g))} + 2^{3 - \deg(M_2(g))}) - (2^{3 - \deg(M_1(g)) - \deg(M_2(g))}) - 1 \\ &= (2 + 2^2) - 2^0 - 1 \\ &= 4. \end{aligned}$$

Therefore, by Corollary 2.6, $Z_{F_2 C_3} = \{g + 1, g^2 + 1, g^2 + g, g^2 + g + 1\}$ is the set of all generators of zero divisor codes over $F_2 C_3$.

Let $W = F_2C_3$. It can be verified that there are two distinct zero divisor codes over F_2C_3 namely $C_1 = W(g + 1) = W(g^2 + 1) = W(g^2 + g) = \{0, g + 1, g^2 + g, g^2 + 1\}$ and $C_2 = W(g^2 + g + 1) = \{0, g^2 + g + 1\}$ with the minimal distance of C_1 and C_2 are 2 and 3 respectively [5]. Hence, in practice C_2 is better than C_1 .

Example 2.11. Consider F_2C_{12} . Recall that $g^3 - 1 = (g^2 + g + 1)(g + 1)$ from Example 2.10. Note that $\gcd(2, 12) = 2$ and $12 = 4 \times 3$. Then by Theorem 2.4(ii), $g^{12} - 1 = (g^2 + g + 1)^4(g + 1)^4$. By Corollary 2.9, $|Z_{F_2C_{12}}| = 2^{(4-1) \times 3} (|Z_{F_2C_3}| + 1) - 1 = 2^9(4 + 1) - 1 = 2559$. Note that

$$\Upsilon_1 = \left\{ b(g)(g^2 + g + 1) \mid b(g) = \sum_{k=0}^9 \epsilon_k g^k, \epsilon_k \in F_2 \right\} \text{ and}$$

$$\Upsilon_2 = \left\{ b(g)(g + 1) \mid b(g) = \sum_{k=0}^{10} \epsilon_k g^k, \epsilon_k \in F_2 \right\}.$$

Hence, similar to Example 2.10, $Z_{F_2C_{12}} = \chi_1 \cup \chi_2$ is the set of all generators of zero divisor codes over F_2C_{12} where $\chi_1 \cap \chi_2 = \emptyset$ and

$$\chi_1 = \Upsilon_1 - \{0\} = \left\{ b(g)(g^2 + g + 1) \mid b(g) = \sum_{k=0}^9 \epsilon_k g^k, \epsilon_k \in F_2, b(g) \neq 0 \right\};$$

$$\chi_2 = \Upsilon_2 - \Upsilon_1$$

$$= \left\{ b(g)(g + 1) \mid b(g) = \sum_{k=0}^{10} \epsilon_k g^k, \epsilon_k \in F_2, b(g) \neq \left(\sum_{l=0}^8 \delta_l g^l \right) (g^2 + g + 1), \delta_l \in F_2 \right\}$$

Then $|\chi_1| = 2^{10} - 1 = 1023$ and $|\chi_2| = 2^{11} - 2^9 = 1536$. Note that $|\chi_1| + |\chi_2| = 1023 + 1536 = 2559$ is consistent with $|Z_{F_2C_{12}}|$ obtained using Corollary 2.9.

3 Identification of Check Elements For Zero Divisor Codes Over F_qC_m

Throughout this section, let $u \in F_qC_m$ be a zero divisor with $d(g) = \gcd(u(g), g^m - 1)$ and $C = Wu$ is its corresponding zero divisor code. Then by Theorem 2.2, $d(g) \neq 1$. The following result shows that u is a zero divisor implies d is also a zero divisor.

Proposition 3.1. *Let $u \in F_q C_m$ such that u is a zero divisor and $d(g) = \gcd(u(g), g^m - 1)$. Then d is a zero divisor in $F_q C_m$.*

Proof. Note that $\deg(d(g)) \leq \deg(u(g))$ as $d(g) = \gcd(u(g), g^m - 1)$. It is then clear that $\gcd(d(g), g^m - 1) = d(g) \neq 1$. Thus by Theorem 2.2, d is a zero divisor in $F_q C_m$. \square

Next, the remaining discussion is proceed toward developing a method to explicitly identify a check element for C . Note that as $d(g)$ is a divisor of both $g^m - 1$ and $u(g)$, there exist $d'(g), d''(g) \in F_q[g]$ such that

$$d(g)d'(g) = g^m - 1 \quad (3.1)$$

and

$$d(g)d''(g) = u(g). \quad (3.2)$$

By comparing the degrees of polynomials on the both sides of the Equation 3.1 and Equation 3.2 respectively, it is clear that $d', d'' \in F_q C_m$.

Note that d' is a zero divisor partner of d as Equation (3.1) in $F_q[g]$ implies $dd' = 0$ in $F_q C_m$. In addition, d' is a principal zero divisor partner of d by Equation (3.1) [3]. On the other hand, by Equation (3.2), $\gcd(u(g), g^m - 1) = \gcd(d(g)d''(g), g^m - 1) = d(g)$ which implies that $d''(g)$ and $g^m - 1$ are coprime. Then by Theorem 2.2, d'' is a unit. Therefore, we have proven the following result:

Lemma 3.2. *Let $u \in F_q C_m$ such that u is a zero divisor and $d(g) = \gcd(u(g), g^m - 1)$.*

- (i) *If $d(g)d'(g) = g^m - 1$, then $d' \in F_q C_m$ is a principal zero divisor partner of d ;*
- (ii) *If $d(g)d''(g) = u(g)$, then $d'' \in F_q C_m$ is a unit.*

Our next result shows that regardless of whether $u(g)$ divides $g^m - 1$, both u and d share the same set of principal zero divisor partners.

Theorem 3.3. *Let $u \in F_q C_m$ such that u is a zero divisor and $d(g) = \gcd(u(g), g^m - 1)$. Then $d' \in F_q C_m$ is a check element of $C = Wu$ if and only if d' is a principal zero divisor partner of d .*

Proof. Let $d'' \in F_q C_m$ such that $u = dd''$. By Lemma 3.2(ii), d'' is a unit. Thus the corresponding group ring matrix $\sigma_{C_m}(d'')$ is also a unit. Hence, we have

$$\text{rank}(u) = \text{rank}(\sigma_{C_m}(u)) = \text{rank}(\sigma_{C_m}(d)\sigma_{C_m}(d'')) = \text{rank}(\sigma_{C_m}(d)) = \text{rank}(d). \quad (3.3)$$

Assume that d' is a principal zero divisor partner of d . Then $dd' = 0$ and $\text{rank}(d) + \text{rank}(d') = |C_m|$. Note that $ud' = dd''d' = (dd')d'' = 0$ and by Equation (3.3), $\text{rank}(u) + \text{rank}(d') = \text{rank}(d) + \text{rank}(d') = |C_m|$. Conversely, assume that d' is a check element of C . Then $ud' = 0$ and $\text{rank}(u) + \text{rank}(d') = |C_m|$. Note that $ud' = dd''d' = (dd')d'' = 0$. Since d'' is a unit, then $dd' = 0$. In addition, by Equation (3.3), $\text{rank}(d) + \text{rank}(d') = \text{rank}(u) + \text{rank}(d') = |C_m|$. \square

Lastly, the following result gives a method to explicitly identify a check element for each zero divisor code over $F_q C_m$.

Theorem 3.4. *Let $u \in F_q C_m$ such that u is a zero divisor and $d(g) = \gcd(u(g), g^m - 1)$. Then $v \in F_q C_m$ is a check element of $C = Wu$ if $d(g)v(g) = g^m - 1$.*

Proof. Suppose that $v \in F_q C_m$ is a zero divisor partner of u such that $d(g)v(g) = g^m - 1$. Then by Lemma 3.2(i), v is a principal zero divisor partner of d and thus v is a check element of C by Theorem 3.3. \square

In short, using Theorem 3.3 and Theorem 3.4, an algebraic algorithm to explicitly identify a check element of a zero divisor code over $F_q C_m$ is designed as follows:

Algorithm 3.5. *Let $C = Wu$ be a zero divisor code over $F_q C_m$. For each group ring element $a = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q C_m$, recall that $a(g) = \sum_{i=0}^{m-1} \alpha_i g^i \in F_q[g]$ represents the polynomial form of a in $F_q[g]$ and vice-versa.*

Step 1: Find $d(g) = \gcd(u(g), g^m - 1)$ using Euclidean Algorithm.

Step 2: Find $d' \in F_q C_m$ such that $d(g)d'(g) = g^m - 1$.

Then d' is a check element of C .

Example 3.6. *Let $C = Wu$ be a zero divisor code over $F_2 C_{12}$ with $u = 1 + g + g^2 + g^5 + g^8$.*

Step 1: By using Euclidean Algorithm, we have

$$\gcd(u(g), g^{12} - 1) = 1 + g + g^2.$$

Step 2: Using long division,

$$(1 + g + g^3 + g^4 + g^6 + g^7 + g^9 + g^{10})(1 + g + g^2) = g^{12} - 1.$$

Thus $v = 1 + g + g^3 + g^4 + g^6 + g^7 + g^9 + g^{10} \in F_2C_{12}$ is a check element for C by Algorithm 3.5. Note that $uv = 0$ with

$$\sigma_{C_{12}}(u) = \begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{pmatrix} \text{ and}$$

$$\sigma_{C_{12}}(v) = \begin{pmatrix} 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}.$$

It can be verified that $\text{rank}(\sigma_{C_{12}}(u)) = 10$ and $\text{rank}(\sigma_{C_{12}}(v)) = 2$. Hence

$$\text{rank}(u) + \text{rank}(v) = \text{rank}(\sigma_{C_{12}}(u)) + \text{rank}(\sigma_{C_{12}}(v)) = 12 = |C_{12}|.$$

Therefore, $v = 1 + g + g^3 + g^4 + g^6 + g^7 + g^9 + g^{10}$ is a check element of $C = W(1 + g + g^2 + g^5 + g^8)$ over F_2C_{12} . Hence, $g^2 + g^5 + g^8 \in F_2C_{12}$ but $g^2 + g^5 + g^8 \notin C$ as

$$(g^2 + g^5 + g^8)(1 + g + g^3 + g^4 + g^6 + g^7 + g^9 + g^{10}) = 1 + g^2 + g^3 + g^5 + g^6 + g^8 + g^9 + g^{11} \neq 0.$$

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